

The “Collatz Problem” of Domain Theory: Is **FS** = **RB**?

Types and Topology:

A Workshop in honor of Martín Escardó's 60th birthday

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Zou-Li-Guo and the frontier $FS = RB$

A personal prologue: Martín's 5% rule

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$$FS = RB ?$$

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- His advice (roughly):
 - Spend 95% of your time finishing your thesis.
 - Spend at most 5% of your time playing with this problem.

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- His advice (roughly):
 - Spend 95% of your time finishing your thesis.
 - Spend at most 5% of your time playing with this problem.
- This talk is my belated 5%:
now that I have earned my tenure, I can afford to 'waste' time *intelligently* on this Collatz-like puzzle in domain theory.

Why a “Collatz problem” for domain theory?

- Like the Collatz conjecture, a.k.a. $3n + 1$ -problem¹:
 - Very easy to state.
 - No consensus that current tools are strong enough.
 - **Dangerous** to one's research time!

¹<https://www.youtube.com/watch?v=5mFpVDpKX70>

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- For **FS = RB?**:
 1. Domain theory may simply *not be ready* yet.
 2. Anyone who devotes serious time to it risks *wasting* that time.
 3. Yet the question is *fundamental*: it compares two core finitary models of continuity in denotational semantics.

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- I will:
 - Recall the key players (**FS**, **RB**, Lawson's disc domain),
 - Explain why Lawson's disc domain is FS with some new insight,
 - And indicate how far we are from resolving **FS = RB**.

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From posets to domains

Definition (Dcpo)

A poset (P, \leq) is *directed complete* if every directed subset has a supremum.

A dcpo which has the least element \perp is a pointed dcpo.²

- directed: $D \neq \emptyset$ and $\forall d_1, d_2 \in D. \exists d_3 \in D. d_1, d_2 \leq d_3$.
- supremum: least upper bound of D (if it exists).

²In this talk, all dcpo's are pointed.

The dcpo of closed discs

Let

$$L := \left\{ D(a, r) := \{x \in \mathbb{R}^2 : d(x, a) \leq r\} \mid a \in \mathbb{R}^2, r \geq 0 \right\} \cup \{\mathbb{R}^2\}$$

be the collection of all *closed discs* in \mathbb{R}^2 , including degenerate ones $\{a\}$ and the super-huge one \mathbb{R}^2 .

Order L by reverse inclusion:

- $\mathbb{R}^2 \supseteq D$ for all $D \in L$ (so \mathbb{R}^2 is bottom);
- $D(a, r) \leq D(b, s)$ iff $D(a, r) \supseteq D(b, s)$.

Then (L, \leq) is a pointed dcpo: directed suprema are given by filtered intersections of discs.

The dcpo of closed discs (geometry)

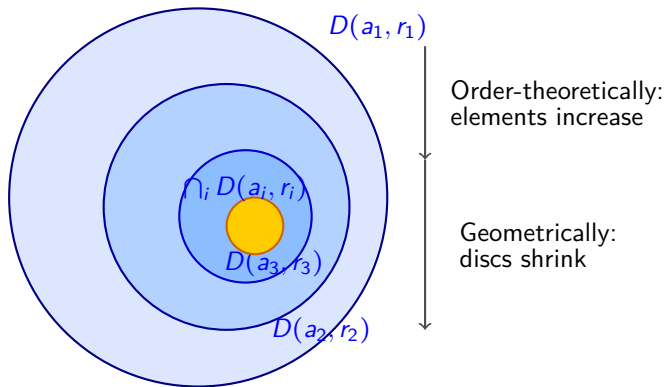


Figure 1: The disc dcpo $(L, \sqsupseteq) : D(a, r) \leq D(b, s) \iff \|a - b\| + s \leq r$.

Domains and the way-below relation

Way-below relation

For a dcpo P :

$$y \ll x \iff \forall D \subseteq_{\text{dir}} P. \left(\bigvee^{\uparrow} D \geq x \implies \exists d \in D. d \geq y \right).$$

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Definition (Domain = Continuous dcpo)

A subset $B \subseteq P$ is a *basis* of P if for every $x \in P$,

$$B_x := \downarrow x \cap B$$

contains a directed subset with supremum x .

A dcpo is called a *continuous dcpo* or *domain* if it has such a basis.

The Lawson's disc domain

In (L, \leq) ,

$$D(a, r) \ll D(b, s) \iff \text{int}(D(a, r)) \supseteq D(b, s),$$

and $\mathbb{R}^2 \ll D$ for all $D \in L$.

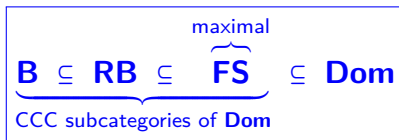
Indeed, for any $D \in L$,

$$D = \bigcap^{\downarrow} \{D' \in L \mid D' \ll D\},$$

which shows that (L, \leq) is a pointed domain, called the *Lawson's disc domain*.

Fact: (L, \leq) is not algebraic: it has not enough compact elements.

RB and FS in the domain zoo



Class	Key Property	Intuitive Idea
Bifinite	Has an approximate identity of idempotent deflations (δ_i) with finite image	$K(P) = \bigcup_i \delta_i[P]$.
RB	Has an approximate identity of deflations (δ_i) with finite image	Retracts of bifinite domains (δ_i not necessarily idempotent)
FS	Has an approximate identity of finitely separating maps δ	Separation by a finite set F_δ
Domain	\ll is approximating	Directed approximations

$\text{AlgFS} = \mathbf{B}$ and $\text{Dom} = \mathbf{RAlg}$

- **Fact:** every continuous domain is a Scott-continuous retract of an algebraic domain:

$$\text{Dom} = \mathbf{RAlg}.$$

AlgFS = B and **Dom = RAlg**

- **Fact:** every continuous domain is a Scott-continuous retract of an algebraic domain:

Dom = RAlg.

- Gierz et al. (Prop. II-2.20):

Theorem

For a dcpo P , the following are equivalent:

1. P is an algebraic FS domain.
2. P is an algebraic domain and has an approximate identity consisting of maps with finite image.
3. P has an approximate identity of idempotent deflations with finite image, i.e., P is bifinite.

So

B = AlgFS.

The open problem (35 years!)

Because

$$\text{Dom} = \text{RAlg}$$

and

$$\text{B} = \text{AlgFS},$$

one asks:

Open Problem (ca. 35 years)

$$\text{Is } \boxed{\text{FS} = \text{RAlgFS} = \text{RB}} ?$$

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$$\text{Is } \boxed{\text{FS} = \text{RAlgFS} = \text{RB}} ?$$

- No counterexample known.
- No general proof known.
- Lawson's disc domain is a natural test case.

Lawson's disc domain as a test case

In Section 4.2.2 of [2], Achim Jung concluded with a concrete example of an FS-domain, due to Jimmie Lawson ([5]).

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It was (and still is) a *likely candidate* for an FS-domain that is not RB.

It is none other than the Lawson disc domain L .

Lawson's approximate identity on L

Let L denote the Lawson disc domain.

For each $\varepsilon > 0$, define $f_\varepsilon : L \rightarrow L$ by

$$f_\varepsilon(D) = \begin{cases} D(a, r + \varepsilon) & \text{if } D = D(a, r) \subseteq B\left(0, \frac{1}{\varepsilon}\right); \\ \mathbb{R}^2 & \text{otherwise.} \end{cases}$$

Why naive definition will not give FS

Consider the functions of the form

$$g_\varepsilon(D) = \begin{cases} D(a, r) & \text{if } D = D(a, r) \subseteq B\left(0, \frac{1}{\varepsilon}\right); \\ \mathbb{R}^2 & \text{otherwise.} \end{cases}$$

(without inflation).

- They also form an approximating identity on L .
- But they **fail** to be finitely separating:
 - We would need a finite separating set F_ε working for uncountably many pairs $D(a, r) \leq D(a, r)$.
 - This is impossible.

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 - This is impossible.
- Moral: to get finite separation, we must use a **robust gap**:

$$D(a, r) \subseteq m_D \subseteq D(a, r + \varepsilon)$$

that persists under small perturbations of the disc.

Robustness via S_ε

Demarcate the region where f_ε inflates discs:

$$S_\varepsilon := \{B(c, r) : \|c\| + r < 1/\varepsilon\}.$$

Thus, S_ε consists of discs lying strictly inside the open ball $B(0, \frac{1}{\varepsilon})$.

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Lemma (Uniform Robustness)

Fix $\varepsilon > 0$. For each $D_p = B(c_0, r_0) \in S_\varepsilon$ define

$$m_{D_p} := B\left(c_0, r_0 + \frac{\varepsilon}{2}\right).$$

Then with $\delta_\varepsilon := \varepsilon/4$ the following holds:

For any $D_p, D_q \in S_\varepsilon$ with parameter triples $p = (c_0, r_0)$ and $q = (c, r)$ satisfying $\|q - p\| < \delta_\varepsilon$ in \mathbb{R}^3 , we have

$$f_\varepsilon(D_q) \leq m_{D_p} \leq D_q.$$

Diagram illustrating the relationship between a set D_q and its image $f_\epsilon(D_q)$ under a function f . The diagram shows three concentric circles. The innermost circle is red and labeled D_q . The middle circle is green and labeled m_{D_p} . The outermost circle is blue and labeled $f_\epsilon(D_q)$. Inside the red circle, there is a point c_0 and a point c , with a double-headed arrow between them. The distance between c_0 and c is labeled $\|q - p\| < \delta_\epsilon$.



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What we already know (Zou–Li–Guo, 2018)

- Each finitely separating $\delta : P \rightarrow P$ can, by the Axiom of Choice, be represented by a function $f_\delta : P \rightarrow P$ with *finite image* such that

$$\forall x \in P. \delta(x) \leq f_\delta(x) \leq x.$$

- If f_δ is additionally *order-preserving*, then δ is called **super finitely separating (SFS)**.
- Theorem.** A dcpo P is an **RB-domain** iff it admits an approximate identity consisting of SFS maps.

The problem FS = RB reduces to bridging FS and SFS.

The current frontier of the $FS = RB$ problem

- After Zou–Li–Guo, the problem is purely **order-theoretic**:

$$FS = RB$$

if and only if

every f.s. map admits a monotone finite-image refinement.

- Known coincidences:

Domain class	FS	RB coincide?
Algebraic	✓	✓
L-domain	✓	✓
(dual) Consistent join-semilattice	✓	✓
General continuous dcpo	✓	?

- In general, no known way to enforce monotonicity without losing finiteness or continuity.

Closing thoughts: a Collatz-style open problem

- We now understand **FS = RB** ϵ -better than 35 years ago:
 - FS-ness of L can be justified through a certain robustness condition.
 - RB is characterised by *super* finitely separating maps.
 - We know many positive coincidence results (algebraic, L-domains, etc.).

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- But the core question

$$FS = RB ?$$

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- Perhaps domain theory is still not ready.
Perhaps we need new tools.

In the meantime, this problem serves as our Collatz: a warning label on how we spend our mathematical time.

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-  J. Lawson, Metric Spaces and FS-domains, *Theoretical Computer Science*, 405(1–2), 73–74, 2008.
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Algebraic domains (backup)

Definition (Compact elements)

An element x of a dcpo P is *compact* if $x \ll x$.

The collection of all compact elements of P is denoted by $K(P)$.

Definition (Algebraic domains)

A dcpo P is an *algebraic domain* if it has $K(P)$ as a basis.

Example & Non-example

1. The powerset lattice $(\mathbb{P}(X), \subseteq)$ is algebraic.
2. The Lawson's disc domain (L, \supseteq) is **not** algebraic.

Continuous = R-Algebraic (backup)

Definition (Scott-continuous retraction)

Let P and Q be dcpo's, and Scott-continuous mappings between them

$$P \begin{array}{c} \xrightarrow{s} \\ \xleftarrow{r} \end{array} Q$$

such that $r \circ s = \text{id}_P$.

We call $r: Q \rightarrow P$ a *Scott-continuous retraction*, and P is a (Scott-continuous) *retract* of Q .

Theorem

Every continuous domain is the retract of an algebraic domain, i.e.,

$$\text{Dom} = \text{RAlg}.$$

Flat natural numbers is bifinite (backup)

The sequence, $(\delta_k(N))_{k \in \mathbb{N}}$, of finite posets has bilimit $N = \mathbb{N}_\perp$.

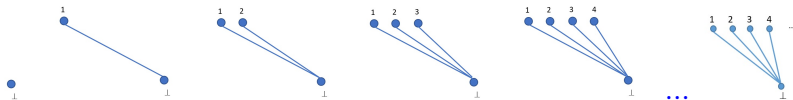


Figure 3: N is bifinite.

A simple FS example on $[0, 1]$ (backup)

Consider the unit interval $[0, 1]$ with the usual order \leq .
For each $\epsilon > 0$, the function

$$f_\epsilon(x) := \max\{x - \epsilon, 0\}$$

is finitely separated from the identity map $\text{id}_{[0,1]}$.
Of course, $([0, 1], \leq)$ being a bc-domain must be FS.

Proof of Uniform Robustness (backup)

Proof.

As usual, $B(c_1, r_1) \subseteq B(c_2, r_2)$ iff $\|c_1 - c_2\| + r_1 \leq r_2$.

If $\|q - p\| < \delta_\varepsilon$ then $\|c - c_0\| < \delta_\varepsilon$ and $|r - r_0| < \delta_\varepsilon$. Thus

$$\|c - c_0\| + r \leq r_0 + 2\delta_\varepsilon = r_0 + 2 \cdot \frac{\varepsilon}{4} = r_0 + \frac{\varepsilon}{2}$$

showing $D_q \subseteq m_{D_p}$.

Since $D_q \in S_\varepsilon$, we have $f_\varepsilon(D_q) = B(c, r + \varepsilon)$. Also

$$\|c - c_0\| + \left(r_0 + \frac{\varepsilon}{2}\right) \leq \delta_\varepsilon + r_0 + \frac{\varepsilon}{2} \leq \left(r + \frac{\varepsilon}{4}\right) + \frac{3\varepsilon}{4} = r + \varepsilon,$$

establishing $m_{D_p} \subseteq f_\varepsilon(D_q)$.

