

Sharp Elements and the Scott Topology of Domains

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*37th Conference on Mathematical Foundations of
Programming Semantics (MFPS)*

Salzburg, Austria (hybrid)

2 September 2021



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BIRMINGHAM

Introduction

Goal

Constructive study of the Scott topology of continuous dcpos.

Relevance of domain theory

Applications in:

- semantics of programming languages;
- higher-type computability;
- topology & locale theory.

Constructivity

- We don't assume **excluded middle** or the **axiom of choice**.
- Our results are valid in any **topos** (e.g. Hyland's effective topos).
- Constructive mathematics has been used for **program extraction**.
- Constructive domain theory may inspire **computable** domain theory.

Example of constructive benefits

- We can prove constructively that the **Scott model** of the programming language **PCF** is sound and **computationally adequate**.
- Given a PCF program t of base type and $n \in \mathbb{N}$,

$$\llbracket t \rrbracket = \llbracket \underline{n} \rrbracket \iff t \text{ reduces to } \underline{n}.$$

- Because the proof is constructive, we get an **interpreter** that runs a program, if we provide it with a proof that the program is *total*.

The proof of totality can use topological/denotational ideas.

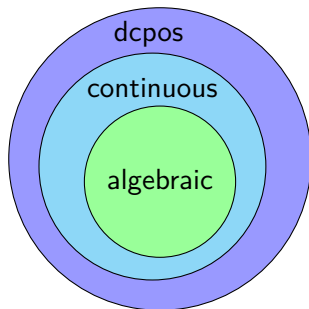
Outline

- 1 Preliminaries: basic domain theory
 - **Examples** of continuous dcpos
 - The Scott topology
- 2 The intrinsic apartness
- 3 Sharp elements
- 4 Strong maximality

Highlight and explain constructive aspects throughout the talk.

Basic objects in domain theory

- In domain theory we study **directed complete posets (dcpos)**: posets with **suprema** for **directed subsets**.
- All our **examples** will be **continuous** dcpos which have a **basis** whose elements can be used to *approximate* any other element.
- In the special case of **algebraic** dcpos every element can be *approximated* using **compact** (“finitary”) elements.



Examples of continuous dcpos I

Example

- The **powerset** $\mathcal{P}(X)$ of any set X , ordered by subset inclusion, is an algebraic dcpo.
- Its compact elements are exactly the *Kuratowski finite* subsets. A subset S is Kuratowski finite if it is **finitely enumerable**, possibly with **repetitions**.
- Since \mathbb{N} has **decidable equality**, we can get rid of repetitions in Kuratowski finite subsets of \mathbb{N} .
- But we **cannot** do this for subsets of **real numbers**, because \mathbb{R} does not have decidable equality.

Examples of continuous dcpos II

Another class of examples is given by considering **rounded ideals** on **abstract bases**.

Example

- Consider the poset (A^*, \preceq) of **finite sequences** on some set A ordered by prefix.
- We get an algebraic dcpo \mathcal{A} whose compact elements are precisely of the form $\downarrow \tau := \{\sigma \in A^* \mid \sigma \preceq \tau\}$ with τ a finite sequence.
- We have an injection from the set of **infinite sequences** to \mathcal{A} given by $\alpha \mapsto \{\sigma \in A^* \mid \sigma \text{ is a prefix of } \alpha\}$.

Example

Taking rounded ideals on $(\mathbb{Q}, <)$ we get a continuous dpcpo of **lower Dedekind reals**.

Examples of continuous dcpos III

Example

- Consider the set $\mathbb{Q} \times_{<} \mathbb{Q} := \{(p, q) \in \mathbb{Q} \times \mathbb{Q} \mid p < q\}$ ordered by $(p, q) \prec (r, s) \iff p < r < s < q$.
- Taking rounded ideals we get a continuous dcpo \mathcal{R} of **partial (two-sided) Dedekind reals** that classically, is isomorphic to the well-known **interval domain**.
- The partial reals are bounded, rounded and transitive, but may fail to be located: a real x is **located** if for every $p, q \in \mathbb{Q}$ with $p < q$ we have $p < x$ or $x < q$.
- We may think of a partial real as a *computation* of rational approximations which may get *stuck*.
- We have an injection from \mathbb{R} to \mathcal{R} given by $x \mapsto \{(p, q) \mid p < x \text{ and } x < q\}$.

The Scott topology

Definition

A subset C of a dcpo D is *Scott closed* if

- it is a *lower set*: if $x \sqsubseteq y \in C$, then $x \in C$;
- it is *closed under directed suprema*: $\bigsqcup S \in C$ for directed $S \subseteq C$.

Definition

A subset U of a dcpo D is *Scott open* if

- it is an *upper set*.
- it is *inaccessible by directed suprema*: if $\bigsqcup S \in U$, then $s \in U$ for some $s \in S$ for every directed $S \subseteq D$.

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- Complements of Scott opens are Scott closed.
- But, complements of Scott closed subsets are Scott open if and only if excluded middle holds.

Introduction to apartness

- For real numbers $x \neq y$ is very **weak**, because it is a *negative* notion.
- Brouwer introduced **apartness** on real numbers, a *positive* notion:

$$x \# y \iff \exists q \in \mathbb{Q} (x < q < y) \vee (y < q < x).$$

- Classically, $x \# y \iff x \neq y$, but constructively, $x \# y$ is stronger.
-
- Apartness on real numbers is used by Brouwer and in Bishop's *Foundations of Constructive Analysis* (1976).
 - Bridges and Vîță developed general topology using apartness relations in *Apartness and Uniformity* (2011).

Intrinsic apartness

Classically, every dcpo with the Scott topology is T_0 -separated:

- 1 if x and y have the same Scott open neighbourhoods, then $x = y$;
- 2 if $x \neq y$, then there is a Scott open neighbourhood of x that does not contain y or vice versa.

- Item 1 holds constructively for **continuous** dcpos.
- Item 2 is equivalent to excluded middle.

Intrinsic apartness

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This leads to our first new notion.

Definition

Two elements $x, y \in D$ are **intrinsically apart**, written $x \# y$, if there is a Scott open that contains one but not the other.

Intrinsic apartness in examples

Example

For two subsets $A, B \in \mathcal{P}(X)$ we have $A \# B$ if and only if $(B \setminus A)$ or $(A \setminus B)$ is inhabited.

Example

- If α and β are two infinite sequences on a set A , then we may say that α and β are apart if $\exists n \in \mathbb{N} (\alpha(n) \neq \beta(n))$.
- Recall the injection $\iota: \alpha \mapsto \{\sigma \in A^* \mid \sigma \text{ is a prefix of } \alpha\}$ into the continuous dcpo \mathcal{A} .
- Then, α and β are apart as sequences if and only if $\iota(\alpha)$ and $\iota(\beta)$ are intrinsically apart.

Example

Similarly, two real numbers are apart in Brouwer's sense if and only if they are intrinsically apart in our sense as partial Dedekind reals.

Cotransitivity and tightness I

Definition

A relation $\#$ is:

- *cotransitive* if $x \# y$ implies $x \# z$ or $z \# y$ for every x, y and z .
- *tight* if $\neg(x \# y)$ implies $x = y$.

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Warning

We only ask that an apartness is irreflexive and symmetric, we do **not** require cotransitivity or tightness.



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Proposition

The intrinsic apartness is cotransitive/tight if and only if excluded middle holds.

But we will get cotransitivity and tightness in special cases discussed later.

Cotransitivity and tightness II

Is the intrinsic apartness at fault for not being cotransitive or tight?
No, because of the following result.

Theorem

If $\#$ is any irreflexive relation on a dcpo with elements $x \sqsubseteq y$ such that $x \# y$, then cotransitivity of $\#$ implies weak excluded middle.

Sharp elements

- For each dcpo, we can define a subset of **sharp** elements.
- Sharp elements occur naturally in the **examples**.
- The intrinsic apartness *is* **tight** and **cotransitive** on sharp elements.

Sharp elements

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Proposition

*An element x of an **algebraic** dcpo is sharp if and only if $c \sqsubseteq x$ is decidable for every compact element c .*

Proposition

Every element is sharp if and only if excluded middle holds.

Examples of sharp elements

Example

The sharp elements of $\mathcal{P}(X)$ are exactly the **decidable** subsets of X .

Example

- We can try to make a **one-sided lower** real into a **two-sided** real: given $L \subseteq \mathbb{Q}$, define $U := \{q \in \mathbb{Q} \mid \exists p \in \mathbb{Q} (p \notin L \wedge p < q)\}$.
- (L, U) is bounded, directed and transitive.
- Classically, (L, U) is **located** too, but:
- (L, U) is located if and only if L is sharp in the dcpo of lower reals.

Slogan

Sharpness generalizes locatedness.

Sharpness is also related to a notion of locatedness in formal topology due to Spitters and Kawai.

Cotransitivity, tightness and sharpness

Theorem

In a continuous dcpo, tightness and cotransitivity hold for sharp elements:

- *If x and y are sharp, then $\neg(x \# y)$ implies $x = y$.*
- *If z is sharp and x and y are arbitrary, then $x \# y$ implies $x \# z$ or $z \# y$.*

Strong maximality, classically

Subspaces of maximal elements

Many topological spaces are homeomorphic to the **maximal** elements of some continuous dcpo with the relative Scott topology.

Example

The real line is homeomorphic to the subspace of maximal partial Dedekind reals.

Strong maximality, classically

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Strong maximality

- In general, the subspace of **maximal** elements may fail to be **Hausdorff**.
- Therefore, Smyth and Heckmann studied **strong maximality**: the subspace of **strongly maximal** elements is both **Hausdorff** and **regular**.
- Strong maximality and maximality coincide if and only if the **Lawson condition** holds.

Strong maximality, constructively

Proposition

Every strongly maximal element is sharp.

Example

- The real line is homeomorphic to the subspace of **strongly** maximal partial Dedekind reals.
- If every maximal partial Dedekind real is strongly maximal, then weak excluded middle follows.

Can we expect (strong maximality = maximality + sharpness)?

Examples: strong maximality, maximality & sharpness

Example

- We have an algebraic dcpo \mathcal{C} built from binary sequences.
- Cantor space $2^{\mathbb{N}}$ is homeomorphic to the subspace of *strongly maximal* elements of \mathcal{C} .
- For \mathcal{C} we have: strong maximality = maximality + sharpness.

Example

- We have an algebraic dcpo \mathcal{B} built from sequences of natural numbers.
- Baire space $\mathbb{N}^{\mathbb{N}}$ is homeomorphic to the subspace of *strongly maximal* elements of \mathcal{B} .
- For \mathcal{B} we have: (strong maximality = maximality + sharpness) \iff *Markov's Principle (MP)*.
- **MP**: Every binary sequence that is not all zeroes must have a one.

Summary

Studied Scott topology on dcpos **constructively** featuring many **examples**.

- **Intrinsic apartness** is a *positive* counterpart to the negation of equality.
- The intrinsic apartness is very well-behaved w.r.t. **sharp elements**.
- **Strongly maximal** elements are sharp and the notion is very useful constructively, often more so than maximality.

Further result in the paper

The **Bridges-Vîță framework** applies to domain theory: in many continuous dcpos, the Scott topology coincides with the **apartness topology** induced by the intrinsic apartness.

Extended version of the MFPS paper



Sharp Elements and the Scott Topology of Continuous Dcpos.
arXiv: [2106.05064](https://arxiv.org/abs/2106.05064) (math.LO).